

Fault-free Hamiltonian Cycles in Pancake Graphs with Conditional Edge Faults

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Outline

- Hamiltonian problems
- Fault tolerant problems
- Pancake graphs
- Problem and previous results
- Main result and proof idea

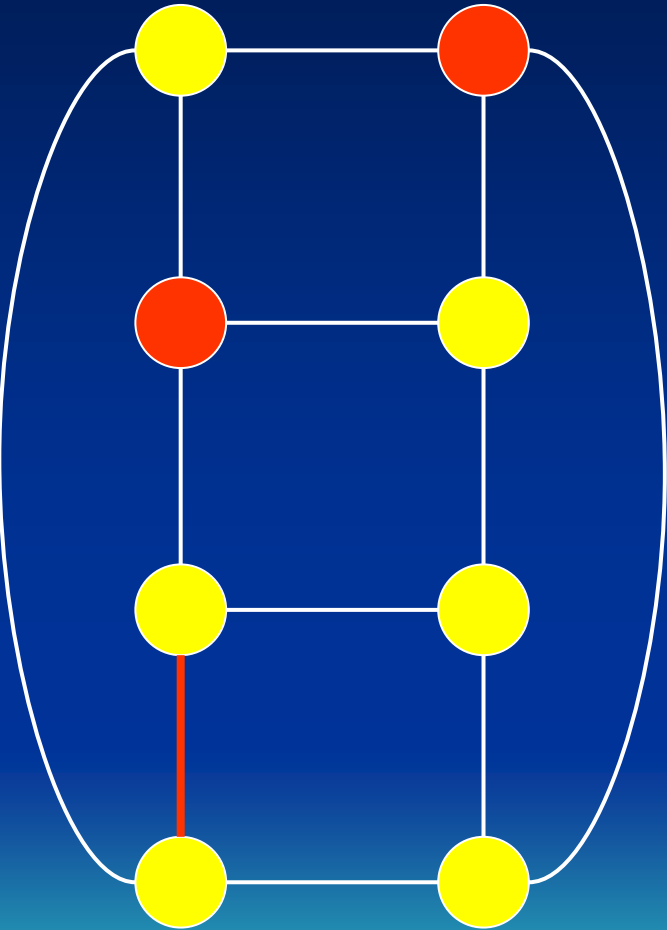


Hamiltonian Problems

- Ring (or linear array) embedding in networks
- Cycle (or path) embedding in graphs
- A cycle (or path) in a graph G is called a *Hamiltonian cycle* (or *Hamiltonian path*) if it contains every vertex of G exactly once.
- G is called *Hamiltonian-connected* if every two vertices of G are connected by a Hamiltonian path.



Fault Tolerant Problems



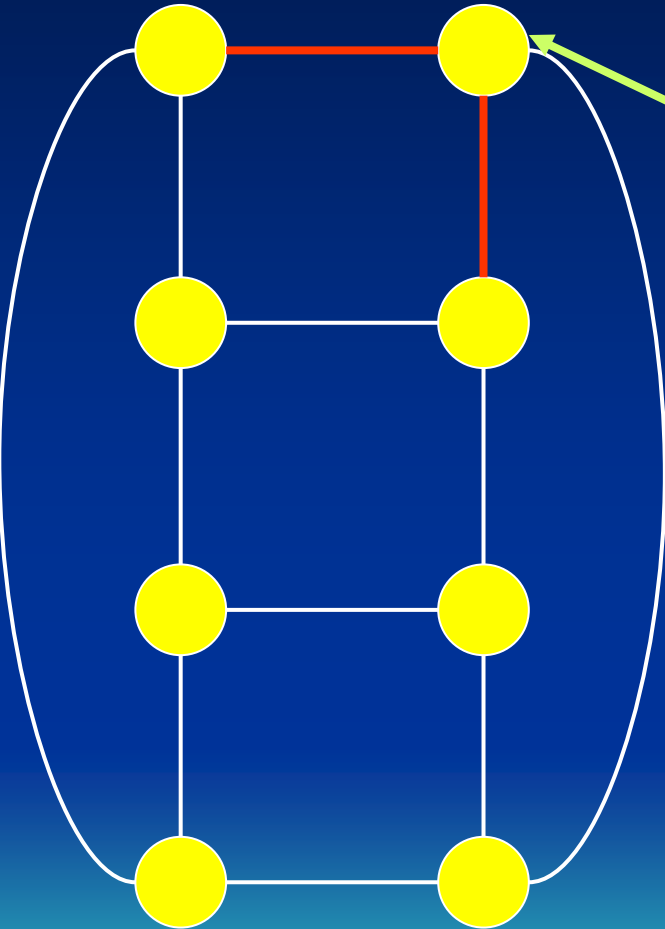
- Vertex faults (node faults).
- Edge faults (link faults).
- After some faults occur, does the network still work?

Fault Models

- *Random fault model*: assumed that the faults might occur everywhere without any restriction.
- *Conditional fault model*: assumed that the distribution of faults must satisfy some properties.
- It is more difficult to solve problems under the conditional fault model than the random fault model.



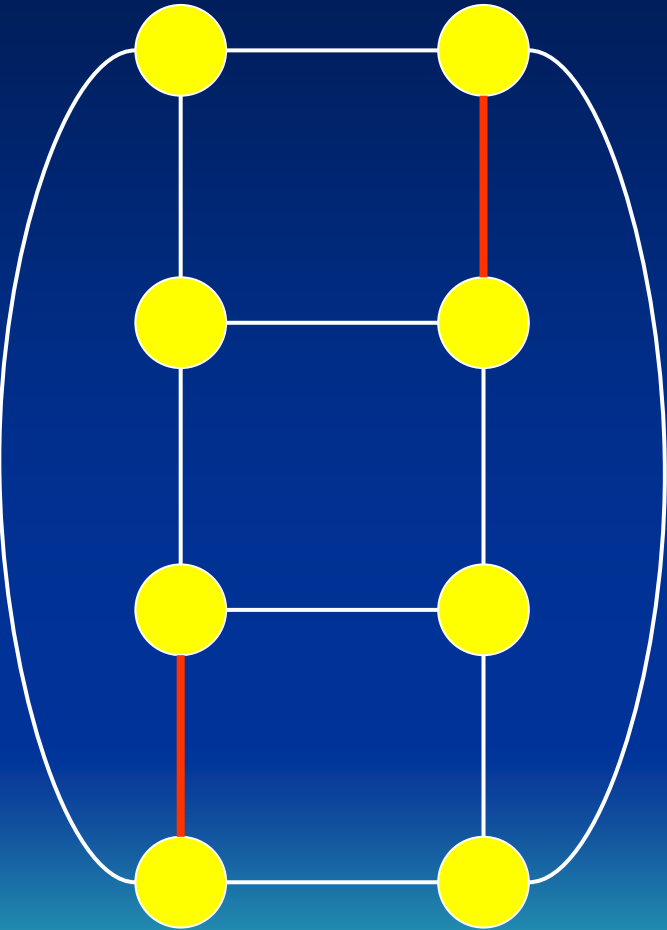
Random Fault – Cycle Embedding



No cycle can pass this vertex

- No Hamiltonian cycle.

Conditional Fault – Cycle Embedding



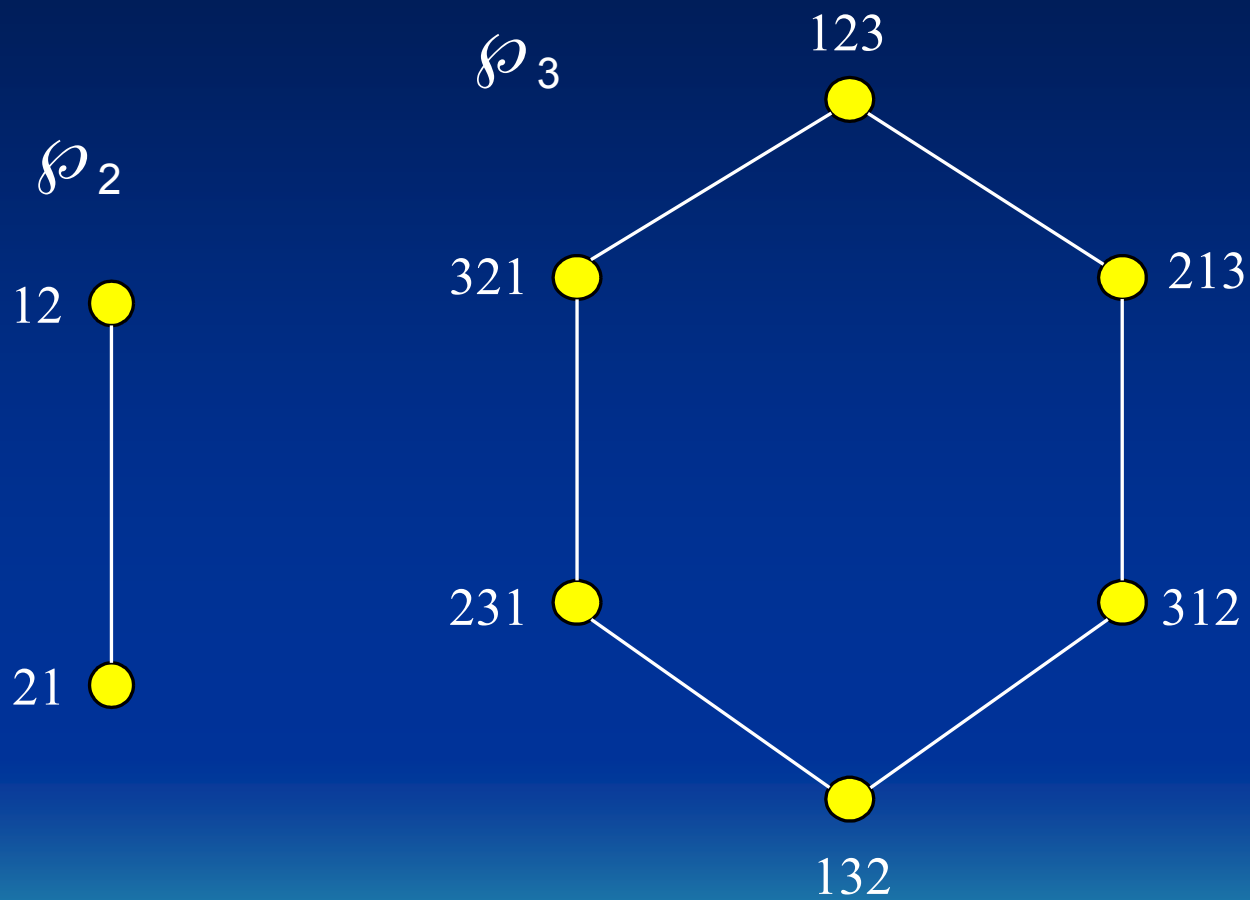
Each vertex has two healthy edges

- No Hamiltonian cycle.

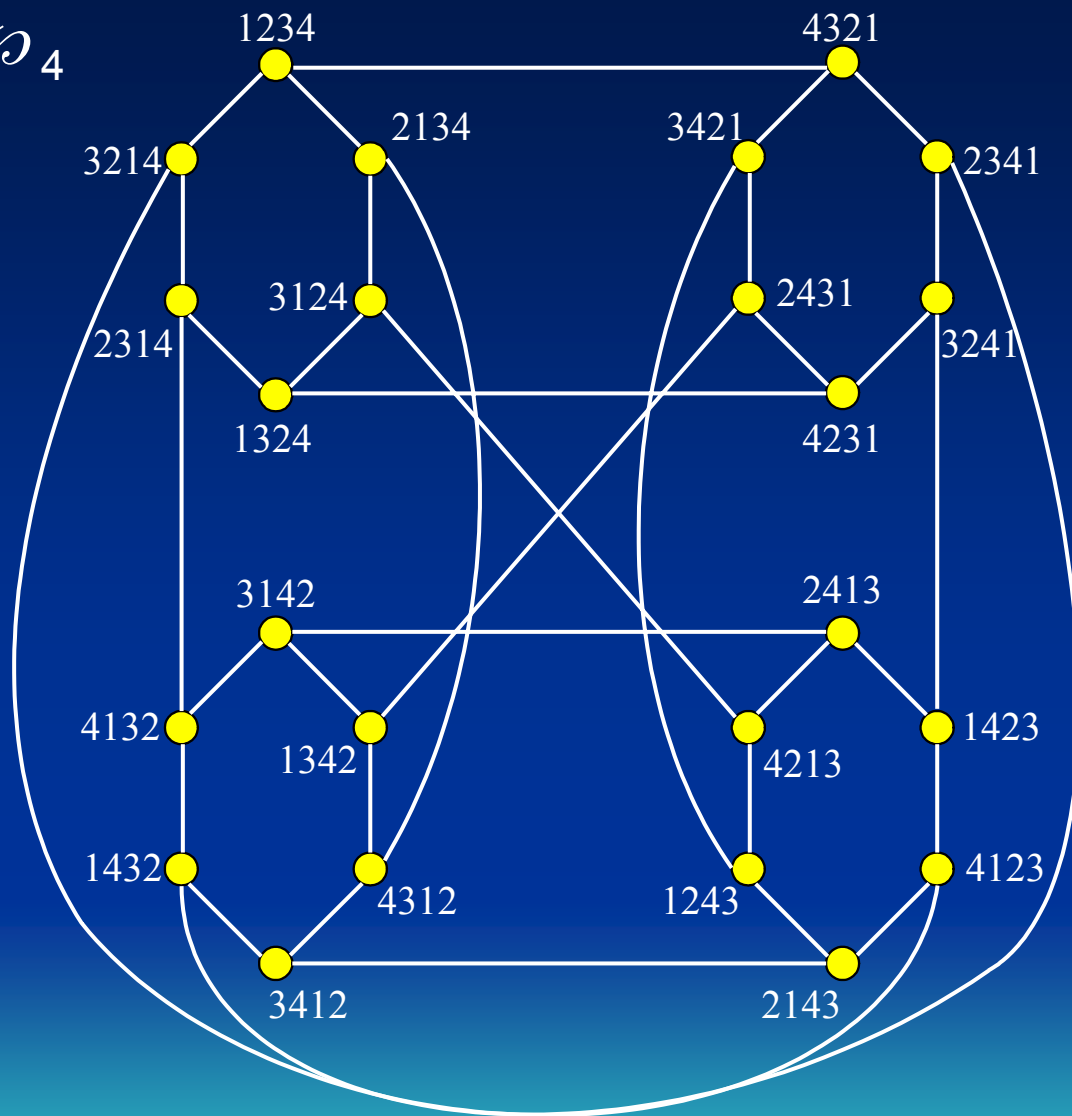
Pancake Graphs

- A n -dimensional pancake graph, denoted by \wp_n , has the vertex set $V(\wp_n) = \{a_1a_2\dots a_n \mid a_1a_2\dots a_n \text{ is a permutation of } 1, 2, \dots, n\}$, and edge set $E(\wp_n) = \{(a_1a_2\dots a_n, b_1b_2\dots b_n) \mid a_1a_2\dots a_k = b_kb_{k-1}\dots b_1 \text{ and } a_{k+1}a_{k+2}\dots a_n = b_{k+1}b_{k+2}\dots b_n \text{ for some } 2 \leq k \leq n\}$.





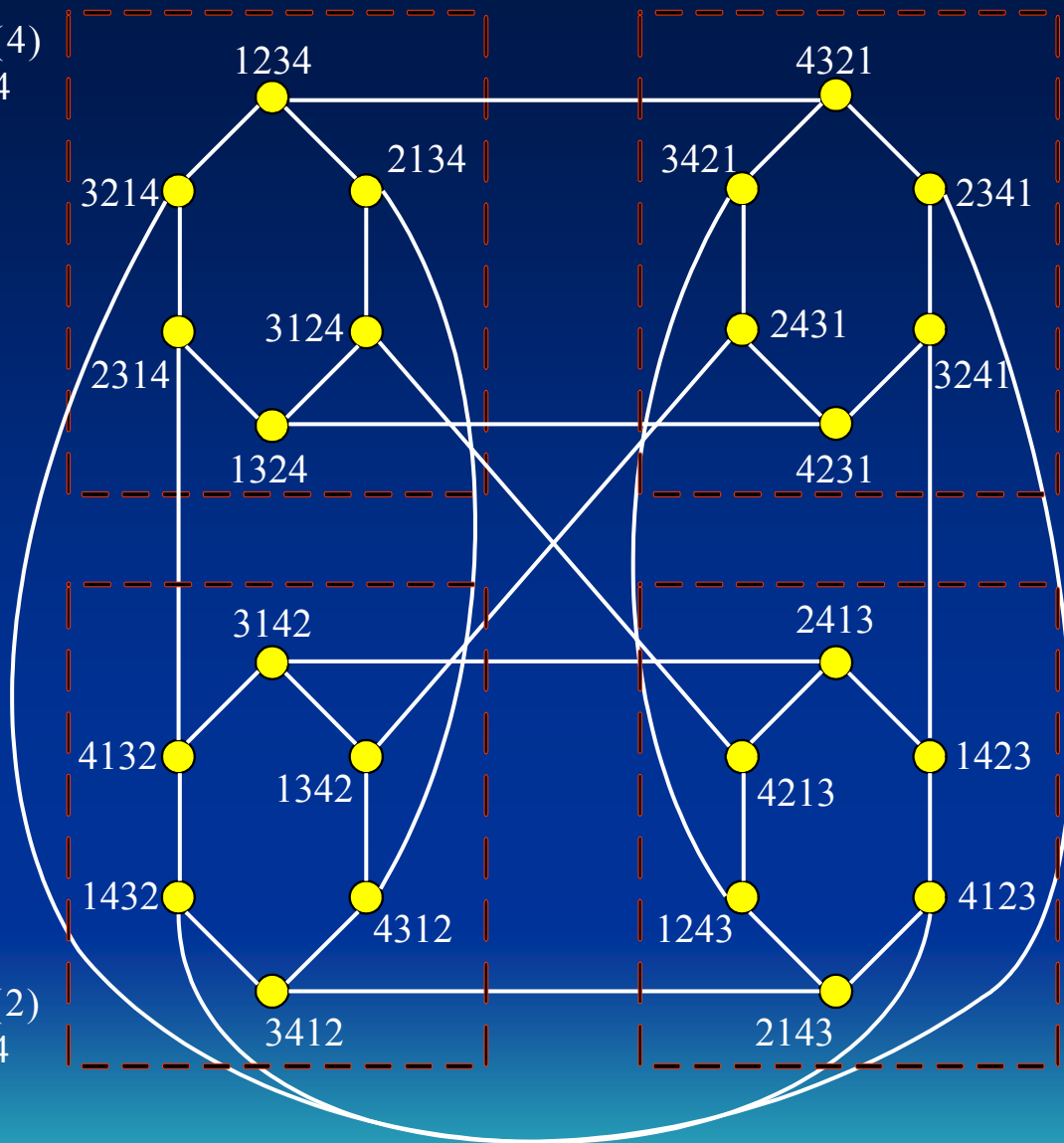
\mathfrak{S}_4



Some Properties of \wp_n

- \wp_n is regular of degree $n-1$.
- \wp_n has $n!$ vertices and $(n-1)n!/2$ edges.
- The girth of \wp_n is 6, where $n \geq 3$.
- \wp_n belongs to the class of Cayley graphs.
- \wp_n is vertex symmetric, but not edge symmetric.
- \wp_n is Hamiltonian-connected.
- \wp_n is a recursive structure.



\mathfrak{P}_4 $\mathfrak{P}_4^{(4)}$ $\mathfrak{P}_4^{(1)}$ $\mathfrak{P}_4^{(2)}$ $\mathfrak{P}_4^{(3)}$ 

A famous problem of \wp_n

- The pancake graph is named from the famous “pancake problem” whose answer is exactly the diameter of the corresponding pancake graph ¹.
- The diameter of \wp_n is bounded above by $3(n + 1)/2$. It is still an open problem to compute the exact diameter of \wp_n ².

- ¹ W. H. Gates and C. H. Papadimitriou, “Bounds for sorting by prefix reversal,” *Discrete Mathematics*, vol. 27, pp. 47-57, 1979.
- ² M. H. Heydari and I. H. Sudborough, “On the diameter of the pancake network,” *Journal of Algorithms*, vol. 25, pp. 67-94, 1997.



Our Problem

- Fault component: edge faults only.
- Fault model: conditional fault model.
- Assumption: each vertex has at least two non-faulty edges.
- How many edge faults can \wp_n tolerate while retaining a fault-free Hamiltonian cycle? ($n \geq 4$)
- This is the first result on the fault tolerance of the pancake graph under the conditional fault model.



Previous results on \wp_n

property \ model	Random fault model	Conditional fault model
k -edge-fault-tolerant Hamiltonian	$k \leq n - 3$ §	$k \leq 2n - 7$
k -edge-fault-tolerant Hamiltonian-connected	$k \leq n - 4$ §	$k \leq n - 4$ *

§ C. N. Hung, H. C. Hsu, K. Y. Liang and L. H. Hsu, “Ring embedding in faulty pancake graphs,” *Information Processing Letters*, vol. 86, pp. 271-275, 2003.

* P. Y. Tsai, “Edge-fault-tolerant path/cycle embedding on some Cayley graphs,” Ph.D. Thesis, National Taiwan University, Taipei, Taiwan, to appear (2008).



Lemmas

- **Lemma 1:** $|\tilde{E}_{p,q}(\wp_n)| = (n-2)!$ for all $p, q \in \{1, 2, \dots, n\}$ and $p \neq q$, where $n \geq 3$.
- **Lemma 2:** $\wp_n - F$ is Hamiltonian if $|F| \leq n-3$, and Hamiltonian-connected if $|F| \leq n-4$, where $n \geq 4$ (F denotes a set of edge faults in \wp_n).



Lemmas

- **Lemma 3:** Suppose that $u, v \in V(\wp_n)$ and $\langle u \rangle_n \neq \langle v \rangle_n$, where $n \geq 5$. For any $I \subseteq \{1, 2, \dots, n\}$ and $|I| \geq 2$, there exists a Hamiltonian path from u to v in $\wp_n^I - F$ provided the following two conditions hold:
(C1) $|\tilde{E}_{i,j}(\wp_n) - F| \geq 3$ for all $i, j \in I$ and $i \neq j$;
(C2) $\wp_n^{(r)} - F$ is Hamiltonian-connected for all $r \in I$.
- **Lemma 4:** Suppose that $u, v \in V(\wp_n^{(r)})$ and $u \neq v$, where $r \in \{1, 2, \dots, n\}$ and $n \geq 4$. If $d_{u,v} \leq 2$, then $\langle N^{(n)}(u) \rangle_n \neq \langle N^{(n)}(v) \rangle_n$, where $d_{u,v}$ is the distance between u and v .



Lemmas

- **Lemma 5:** Suppose that $e_1, e_2 \in E(\wp_4)$ and $e_1 \neq e_2$. There exists a Hamiltonian cycle in $\wp_4 - \{e_2\}$ that contains e_1 .
- **Lemma 6:** Suppose that $s, t \in V(\wp_n)$, $s \neq t$, and $\langle s \rangle_1 = \langle t \rangle_1$, where $n \geq 4$. For every $(x, y) \in E(\wp_n)$ with $\{x, y\} \cap \{s, t\} = \emptyset$, there exists a Hamiltonian path from s to t in \wp_n that contains (x, y) .



Main Result

- **Theorem:** $\wp_n - F$ is Hamiltonian provided $|F| \leq 2n - 7$ and $\delta(\wp_n - F) \geq 2$, where $n \geq 4$ (F denotes a set of edge faults in \wp_n).



Proof idea

- The theorem holds for \wp_4 , which is assured by Lemma 2 ($2n - 7 = n - 3$ as $n = 4$).
- Prove by induction on n .
- Suppose the theorem holds for \wp_k , now we construct a Hamiltonian cycle in $\wp_{k+1} - F$, where $k \geq 4$ and $|F| \leq 2k - 5$.



Proof idea

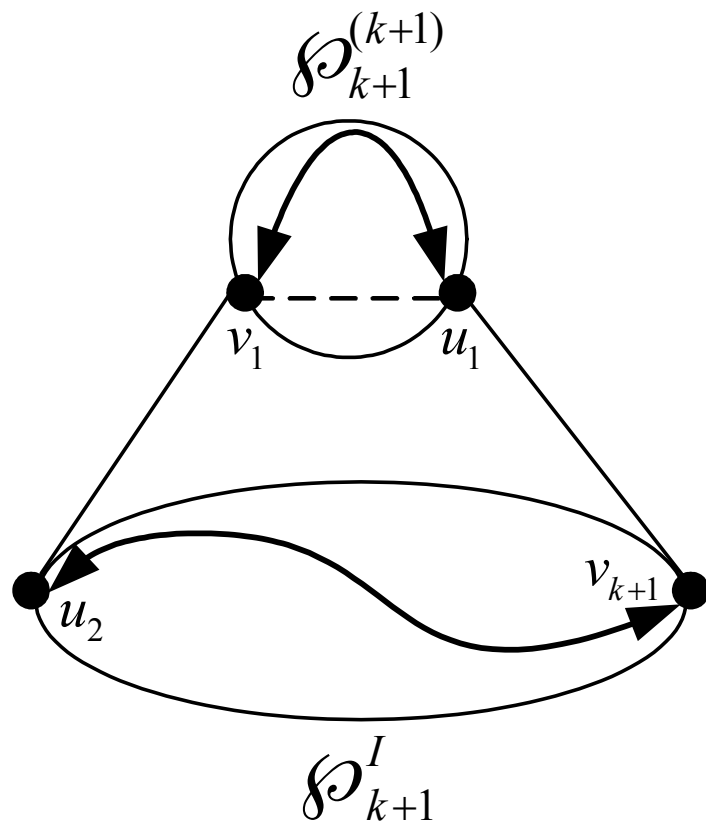
- Assume that $|E(\wp_{k+1}^{(k+1)}) \cap F| \geq |E(\wp_{k+1}^{(k)}) \cap F| \geq \dots \geq |E(\wp_{k+1}^{(1)}) \cap F|$.
- By Lemma 1, we have $|\tilde{E}_{p,q}(\wp_{k+1})| = (k-1)! \geq 2k-2 \geq |F|+3$, i.e., $|\tilde{E}_{p,q}(\wp_{k+1}) - F| \geq 3$ for all $p, q \in \{1, 2, \dots, k+1\}$ and $p \neq q$.
- Four cases are discussed.



Case 1

- $|E(\mathcal{P}_{k+1}^{(k+1)}) \cap F| \leq k - 4$.
- The induction hypothesis assures a Hamiltonian cycle C in $\mathcal{P}_{k+1}^{(k+1)} - F$.
- An edge (u_1, v_1) can be determined from C so that there exist $(v_1, u_2), (u_1, v_{k+1}) \in E^{(k+1)}(\mathcal{P}_{k+1}) - F$ with $u_2, v_{k+1} \in V(\mathcal{P}_{k+1}^I)$, where $I = \{1, 2, \dots, k\}$.
- Lemma 2 assures that $\mathcal{P}_{k+1}^{(j)} - F$ is Hamiltonian-connected for all $1 \leq j \leq k$.
- By Lemma 3, a Hamiltonian path in $\mathcal{P}_{k+1}^I - F$ exists.





Case 2

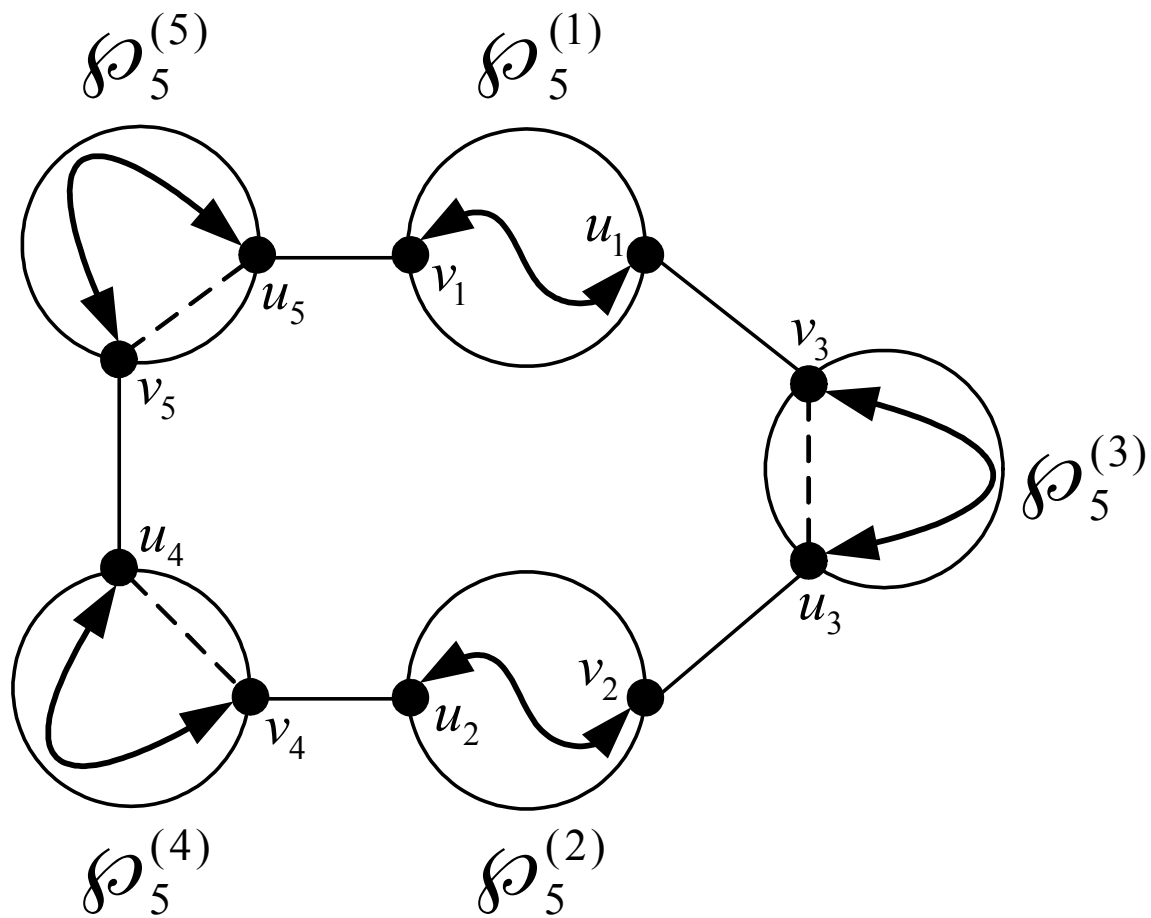
- $k - 3 \leq |E(\wp_{k+1}^{(k+1)}) \cap F| \leq 2k - 7$.
- Subcase 1: $|E(\wp_{k+1}^{(k)}) \cap F| \leq k - 4$.
- If $\delta(\wp_{k+1}^{(k+1)} - F) \geq 2$, the Hamiltonian cycle can be obtained by the construction method of Case 1.
- If $\delta(\wp_{k+1}^{(k+1)} - F) = 1$, we can construct the Hamiltonian cycle by slightly modifying the construction method of Case 1 (use the same figure).



Case 2(continue)

- Subcase 2: $|E(\wp_{k+1}^{(k)}) \cap F| \geq k - 3$.
- We have $|E(\wp_{k+1}^{(k+1)}) \cap F| = k - 3$ or $k - 2$, $|E(\wp_{k+1}^{(k)}) \cap F| = k - 3$ (hence, $\delta(\wp_{k+1}^{(k)} - F) \geq 2$), and $|E^{(k+1)}(\wp_{k+1}) \cap F| \leq 1$.
- First we consider the situation of $|E(\wp_{k+1}^{(k-1)}) \cap F| = k - 3$, which occurs only when $k = 4$.
- We have $|E(\wp_5^{(5)}) \cap F| = |E(\wp_5^{(4)}) \cap F| = |E(\wp_5^{(3)}) \cap F| = 1$, $|E(\wp_5^{(2)}) \cap F| = |E(\wp_5^{(1)}) \cap F| = 0$, and $|E^{(5)}(\wp_5) \cap F| = 0$.
- Assisted by Lemma 5.

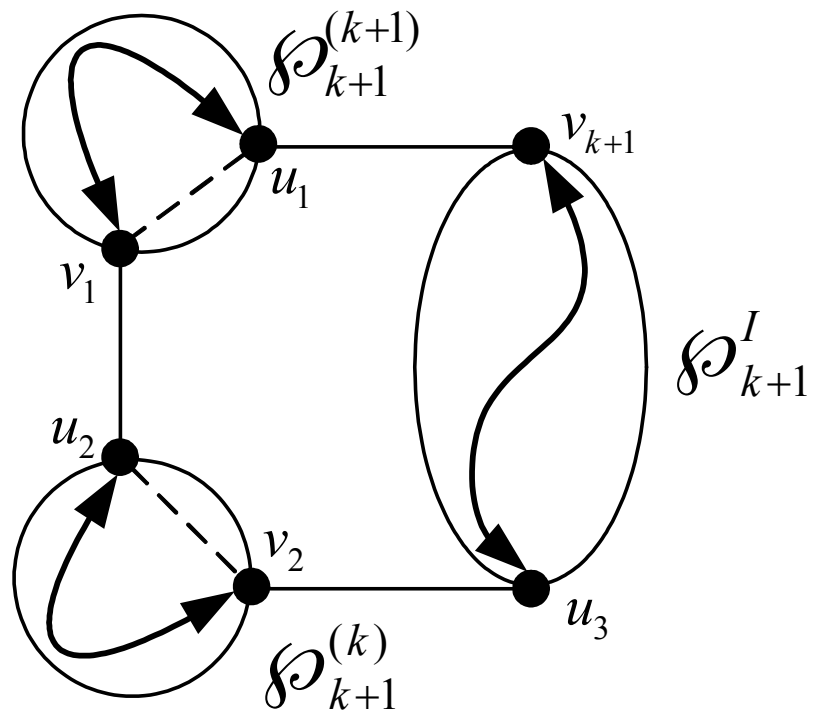


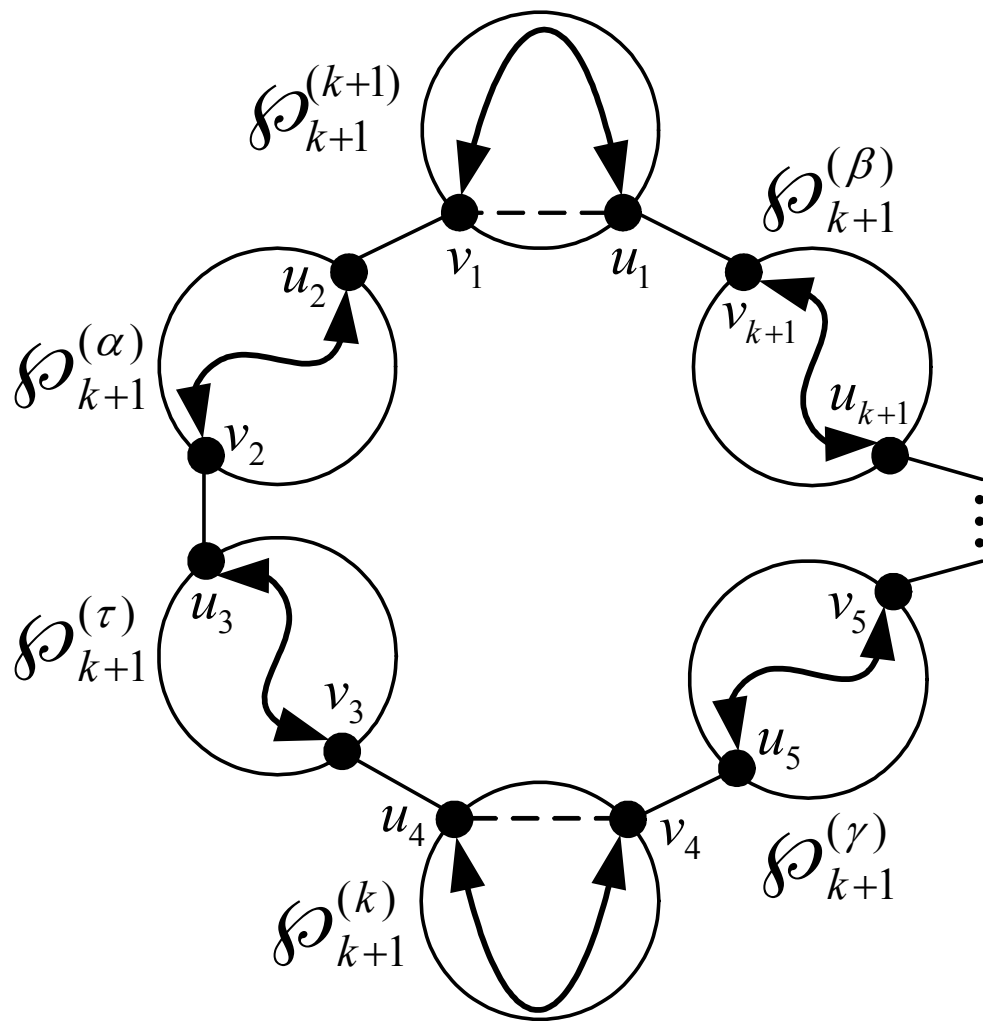


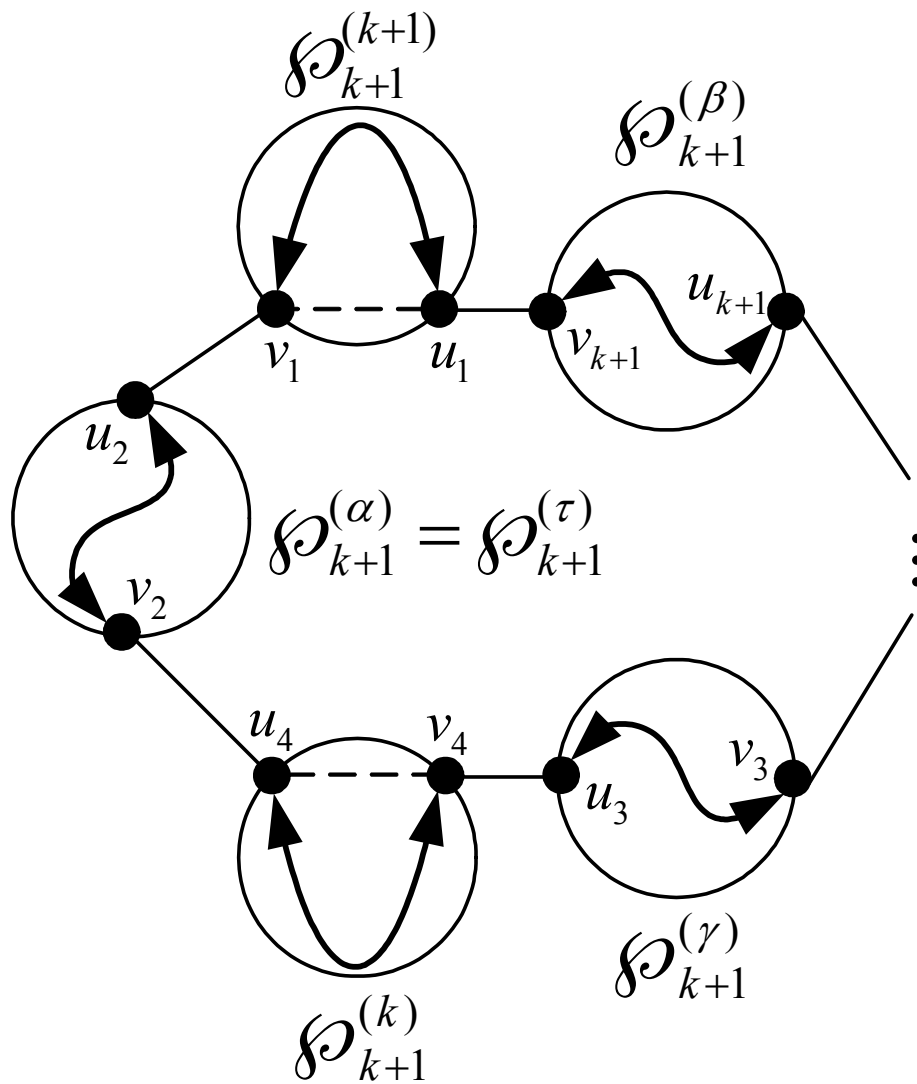
Case 2(continue)

- Now we consider the situation of $|E(\wp_{k+1}^{(k-1)}) \cap F| \leq k-4$.
- When $\delta(\wp_{k+1}^{(k+1)} - F) \geq 2$.
- When $\delta(\wp_{k+1}^{(k+1)} - F) = 1$.









Case 3

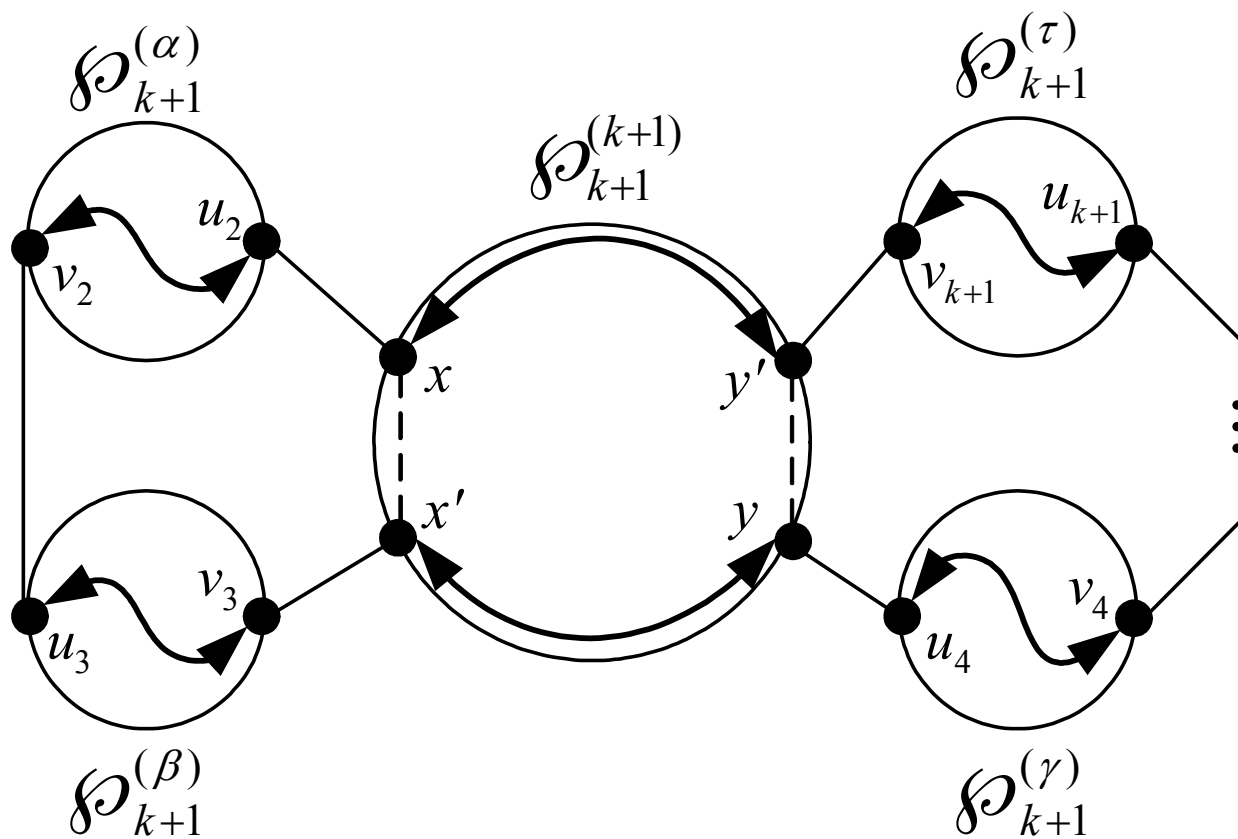
- $|E(\wp_{k+1}^{(k+1)}) \cap F| = 2k - 6.$
- We have $|E^{(k+1)}(\wp_{k+1}) \cap F| \leq 1$ and $|E(\wp_{k+1}^{(j)}) \cap F| \leq 1$ for all $1 \leq j \leq k.$
- When $k \geq 5,$ the Hamiltonian cycle can be obtained by slightly modifying the construction method of Case 1.
- When $k = 4,$ if $|E(\wp_5^{(j)}) \cap F| = 0$ for all $1 \leq j \leq 4,$ the Hamiltonian cycle can be obtained all the same as the situation of $k \geq 5.$ Otherwise, the Hamiltonian cycle can be obtained similar to the construction method of Case 2.

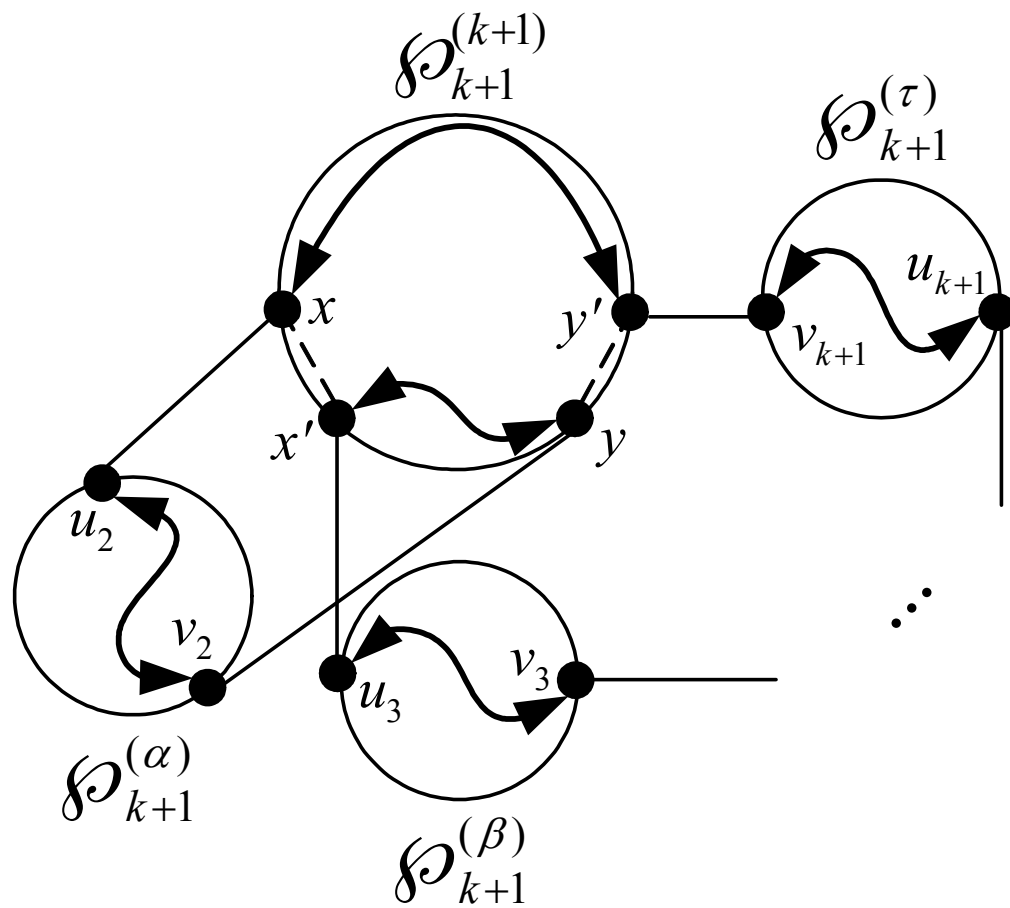


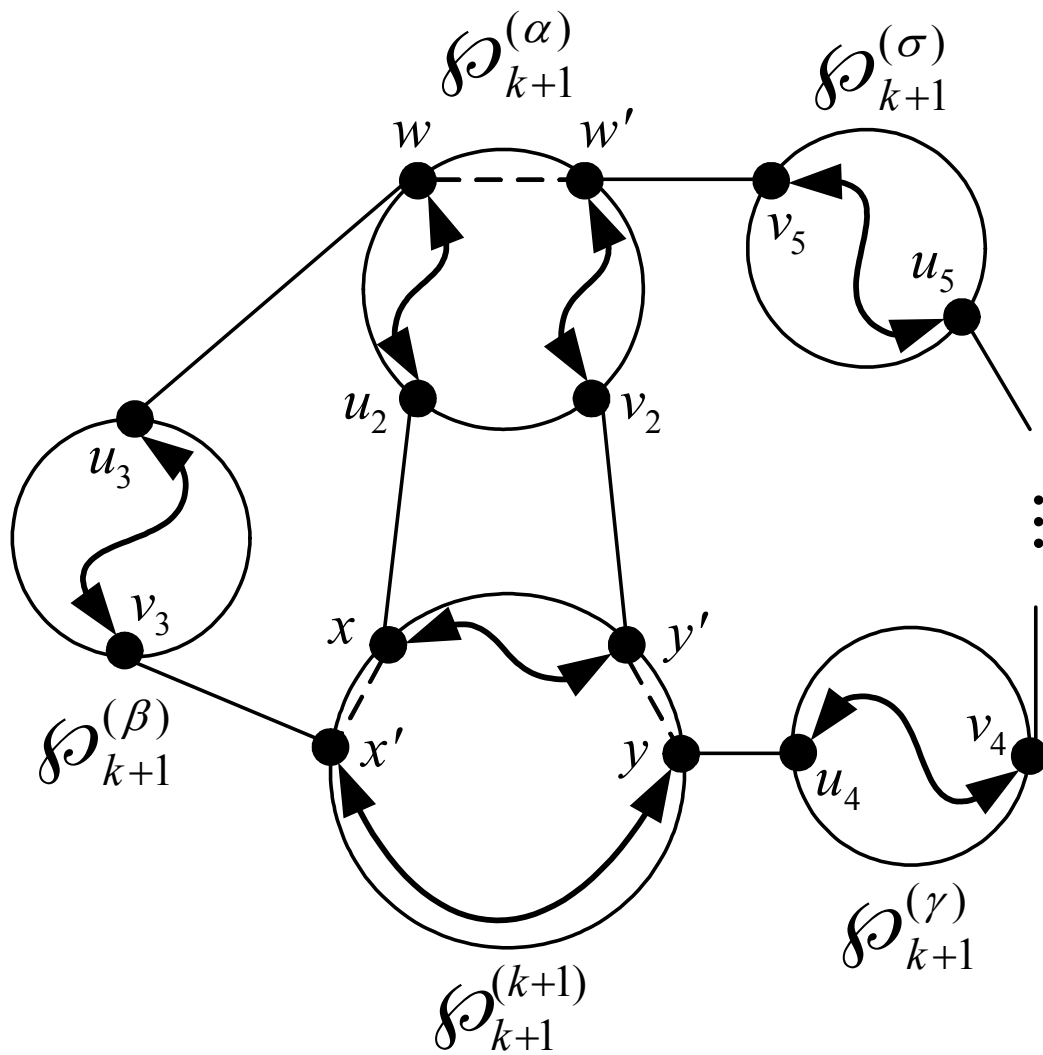
Case 4

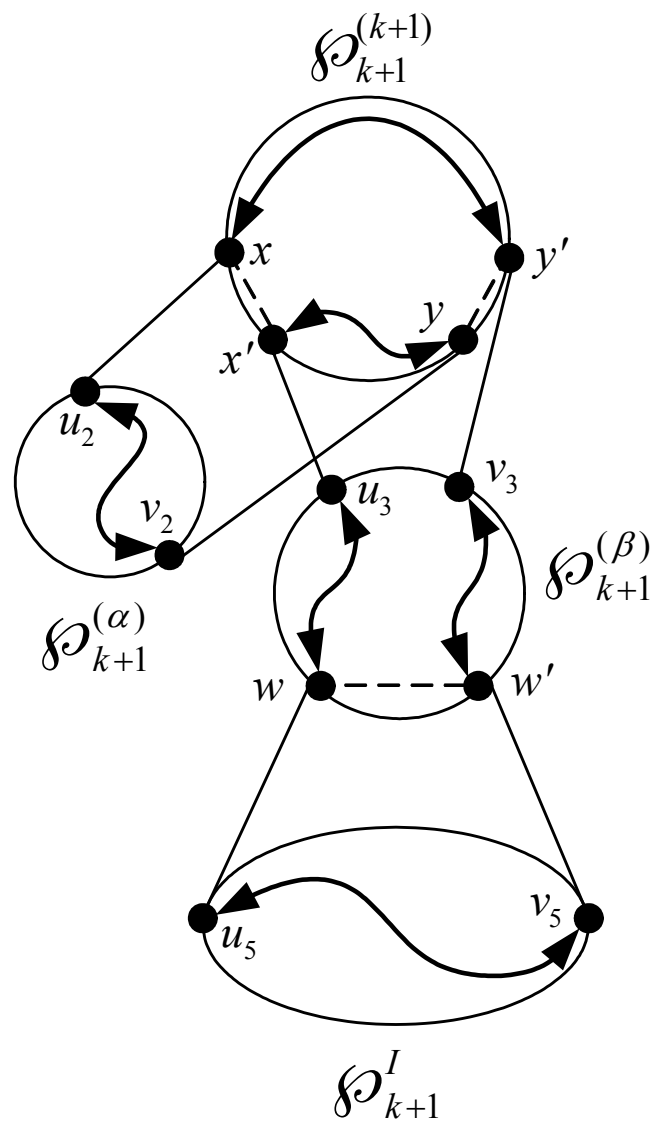
- $|E(\wp_{k+1}^{(k+1)}) \cap F| = 2k - 5$.
- First, two edges $(x, x'), (y, y') \in E(\wp_{k+1}^{(k+1)}) \cap F$ are determined so that $\{x, x'\} \cap \{y, y'\} = \emptyset$ and $\delta(\wp_{k+1}^{(k+1)} - (F - \{(x, x'), (y, y')\})) \geq 2$.
- The induction hypothesis assures a Hamiltonian cycle C in $\wp_{k+1}^{(k+1)} - (F - \{(x, x'), (y, y')\})$.
- If (x, x') or (y, y') is not contained in C , the Hamiltonian cycle can be obtained by the construction method of Case 1.
- Otherwise, five figures are considered.

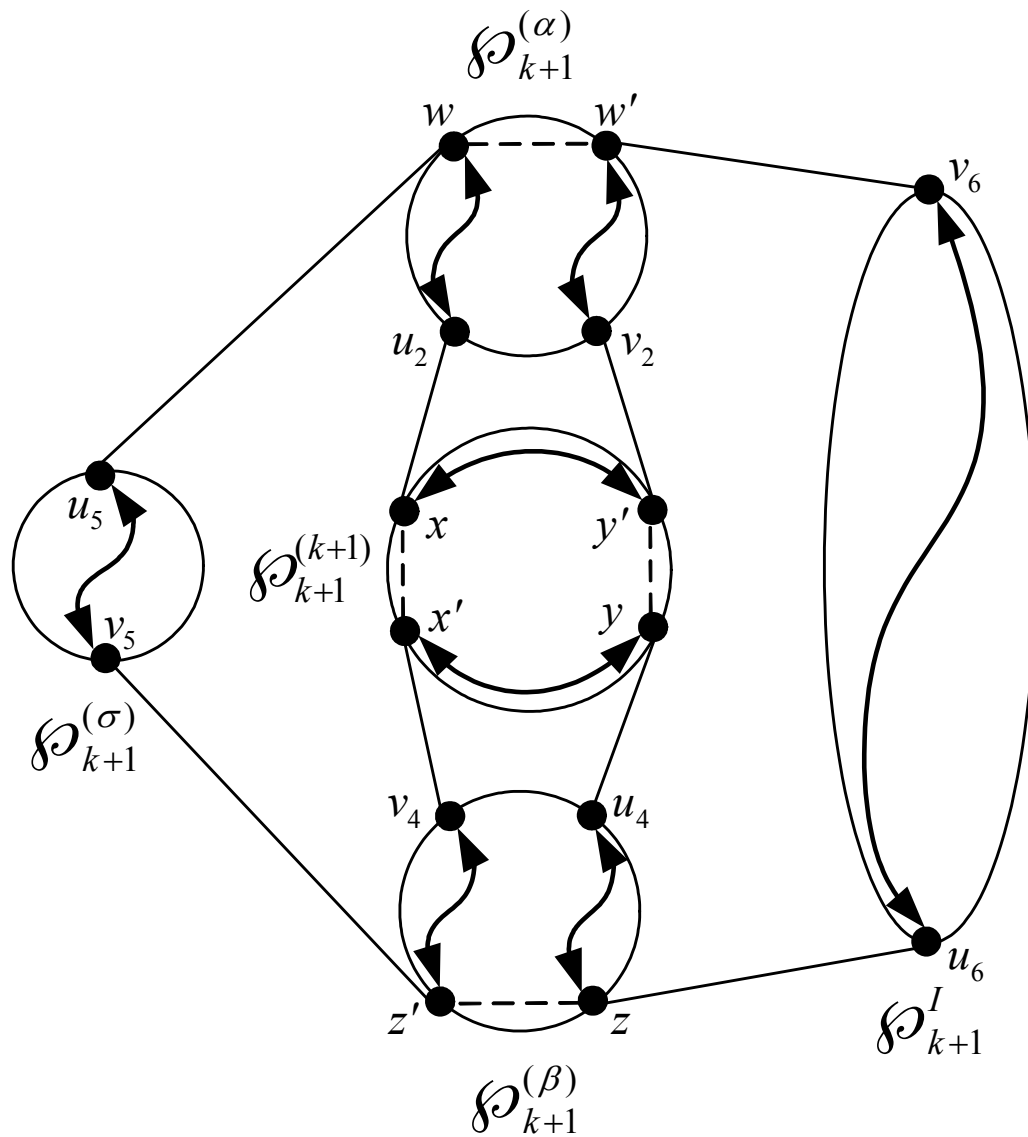




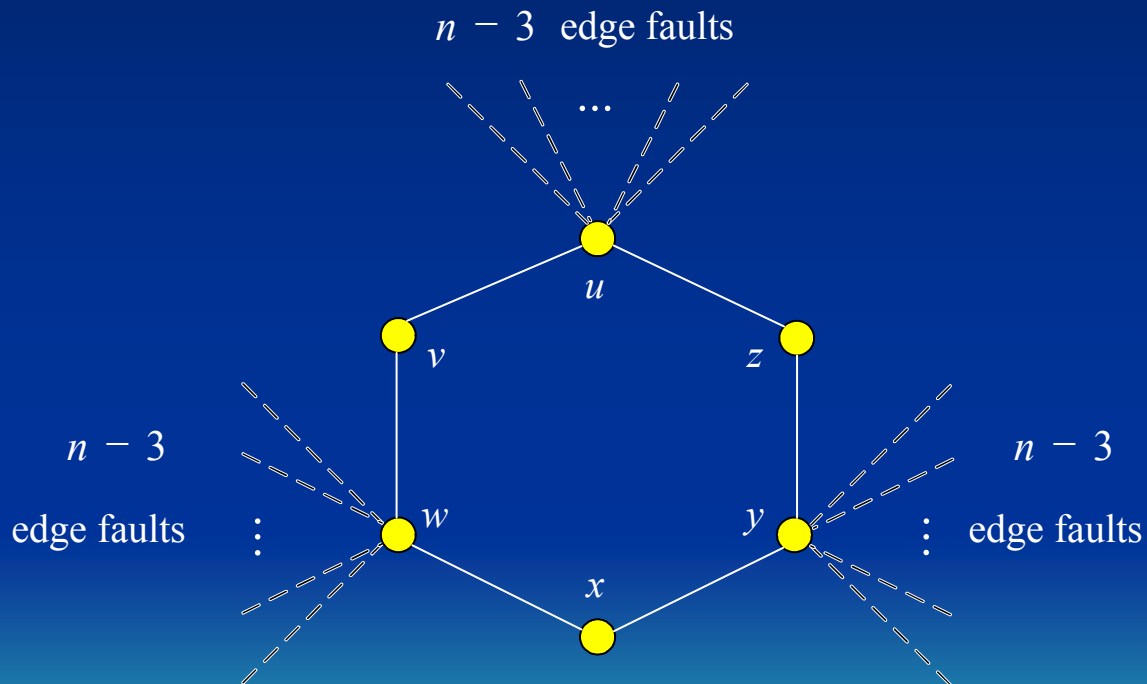








A distribution of $3n - 9$ edge faults over an n -dimensional pancake graph. No fault-free Hamiltonian cycle can be found for this situation.



Open problem

- There is an upper bound of $3n - 10$ on the greatest number of tolerable edge faults for the problem.
- It is an open problem to narrow down the gap between $2n - 7$ and $3n - 10$.

